

Dependency Parsing II

CSC485/2501 Lecture 5

Based on slides by Roger Levy, Yuji Matsumoto, Dragomir Radev, Dan Roth, David Smith and Jason Eisner

Edge-Factored Parsers

- Log-linear models great for n-way classification.
- Also good for predicting sequences.



• Also good for dependency parsing



• Is this a good edge?

yes, lots of green ...



"It was a bright cold day in April and the clocks were striking thirteen."

• Is this a good edge?









• How about this one?

not as good, lots of red ...



Byl jasný studený dubnový den a hodiny odbíjely třináctou was bright cold April day and clock striking thirteen V A A A A N J N V C "It was a bright cold day in April and the clocks were striking thirteen."







- Which edge is better?
 - "bright day" or "bright clocks"?



- Which edge is better?
- Score of an edge: $score(e) = \theta \cdot features(e)$
- Standard algorithms: valid parse with max total score.

jasný studený dubnový den a hodiny odbíjely Byl třináctou bright cold April day and clock striking thirteen was POS: Ν Α V V Α Α Ν С Stem: stud dubn den а hodi odbí třin byl jasn "It was a bright cold day in April and the clocks were striking thirteen."

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- Which edge is better?
- Score of an edge: $score(e) = \theta \cdot features(e)$
- Standard algorithms: valid parse with max total score.





can't have both (no crossing links)

The "projectivity" restriction. Do we really want it?



That glory may-know my going-gray (i.e., it shall last till I go gray) Frequent non-projectivity in Latin, etc.

Non-Projective Parsing Algorithms

Complexity considerations:

- Projective (Proj)
- Non-projective (NonP)

Problem/Algorithm	Proj	NonP
Complete grammar parsing [Gaifman 1965, Neuhaus and Bröker 1997]	Р	NP hard
Deterministic parsing [Nivre 2003, Covington 2001]	O(n)	$O(n^2)$
First order spanning tree [McDonald et al. 2005b]	<i>O</i> (<i>n</i> ³)	$O(n^2)$
Nth order spanning tree (N $>$ 1) [McDonald and Pereira 2006]	Ρ	NP hard

McDonald's Approach (non-projective)

- Example: "Jonh saw Mary "
- We can use the **Chu-Liu-Edmonds** (CLE) algorithm to find the maximum-weight spanning tree.
 - Can be non-projective!
- Can be found in time O(n²).





The **maximum-weight spanning tree**. The best parse with the <mark>highest score</mark>!

Chu-Liu-Edmonds - Contracting Stage

- For each non-ROOT node v, set bestInEdge[v] to be its highest scoring incoming edge.
- If a cycle *C* is formed:
 - contract the nodes in C into a new node v_c
 - edges outgoing from any node in C now get source v_c
 - edges incoming to any node in C now get destination v_c
 - For each node u in C, and for each edge e incoming to u from outside of C:
 - add to *e*.kicksOut the edge bestInEdge[*u*], and
 - set *e*.score to be *e*.score *e*.kicksOut.score.
- Repeat until every non-ROOT node has an incoming edge and no cycles are formed.

- For each non-ROOT node v, set **bestInEdge**[v] to be its ٠ highest scoring incoming edge.
- If a cycle *C* is formed: •
 - contract the nodes in *C* into a new node v_c ٠
 - edges outgoing from any node in C now get source v_c ۰
 - edges incoming to any node in C now get destination v_{c} •
 - For each node *u* in *C* , and for each edge *e* incoming to *u* • from outside of *C*:
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	bestInEdge
/1	
/2	
/3	

	kicksOut
а	
b	
с	
d	
е	
f	
g	
h	
i	

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	bestInEdge
V1	g
V2	
V3	

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oming edge and no		bestInEdg
	V1	
(ROOT)	V2	
	V3	
a:5 b:1 c:1		
		kicksOut
\sim	а	
$\begin{pmatrix} \vee 1 \end{pmatrix} \longrightarrow d : 11 \longrightarrow (\vee 2) \longrightarrow f : 5 \longrightarrow (\vee 3)$	b	
	С	
	d	
e : 4	e	
	f	
h:9	g	
	h	

	bestInEdge
1	g
2	d
3	

0		
	<u>ົ</u>	2

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g

d

g d

g d

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bestInEdge

kicksOut

g

d

g d

g d

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g

d

g

d

g

d

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g

d

h

g

d

g

d

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		bestInEdge
1	V1	g
	V2	d
	V3	f
	V4	h
1	V5	
		kicksOut
	а	g, h
	b	d, h
	С	f
	d	
	е	
	f	
	g	

g

d

h

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		bestInEdge	
,	V1	g	
,	V2	d	
,	V3	f	
,	V4	h	
,	V5		
		kicksOut	
	а	g, h	
	b	d, h	
	с	f	
	d		
	е	f	
	f		
	g		

g d

h

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		bestInEdge	
1	V1	g	-
1	V2	d	
,	V3	f	F
1	V4	h	Ì
	V5	a	Ì
		kicksOut	
	а	g, h	
	b	d, h	
	с	f	
	d		
	е	f	
	f		
	g		

g d

h

Chu-Liu-Edmonds - Expanding Stage

- After the contracting stage, every contracted node will have exactly one **bestInEdge**. This edge will kick out one edge inside the contracted node, breaking the cycle.
 - Go through each bestInEdge of e in the reverse order that we added them
 - lock down e and remove every edge in kicksOut(e) from bestInEdge.

- Go through each **bestInEdge** of *e* in the reverse order that we added them
- lock down e and remove every edge in kicksOut(e) from bestInEdge.



		bestInEdge	•
1	V1	g	5
1	V2	d	
	V3	1	F
1	V4	h	۱
	V5	a	
		kicksOut	
	а	g, h	
	b	d, h	
	с	f	
	d		
	e	f	
	f		
	g		
	h	g	
	i	d	

- Go through each bestInEdge of *e* in the reverse order that we added them
- lock down *e* and remove every edge in kicksOut(*e*) from bestInEdge.



		bestInEdge	•
1	V1	a g	5
	V2	d	
V3		f	
V4		a h	
	V5	a	
		kicksOut	
		MICHOUU	
	а	g, h	
	b	d, h	
	с	f	
	d		
	е	f	
	f		
	g		
	h	g	

d

- Go through each bestInEdge of *e* in the reverse order that we added them
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- Go through each bestInEdge of *e* in the reverse order that we added them
- lock down *e* and <mark>remove</mark> every edge in kicksOut(e) from bestInEdge.



a g d

a h

f

g

d

а

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- lock down *e* and remove every edge in kicksOut(*e*) from bestInEdge.



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Evaluation

• For a sequence labelling task (e.g., PoS tagging), evaluation is straightforward.

Correct

Model

• Exact match (EM) accuracy.

1.	Unionized	VBN	ככ
2.	workers	NNS	NNS
3.	are	VBP	VBP
4.	usually	RB	RB
5.	better	RBR	RBR
6.	paid	VBN	VBN
7.	than	IN	IN
8.	their	PRP\$	PRP\$
9.	non-union	ככ	ככ
10.	counterparts	NNS	NNS

EM Acc = Correct / Total = 90%

Evaluation

• What about dependency parses?



• Solution: count # words are correctly connected to their heads.

Evaluation: Unlabeled dependency accuracy (UAS)



		(a) Reference	(b) System
1	book	0	0
2	me	1	4
3	the	4	4
4	flights	1	1
5	through	6	6
6	Houston	4	4

Accuracy = 5/6 = 83.3%

Evaluation: Labeled dependency accuracy (LAS)



		(a) Reference		(b) System	
		head	label	head	label
1	book	0	root	0	root
2	me	1	iobj	4	nsubj
3	the	4	det	4	det
4	flights	1	obj	1	хсотр
5	through	6	case	6	case
6	Houston	4	nmod	4	nmod

Accuracy = 4/6 = 66.7%